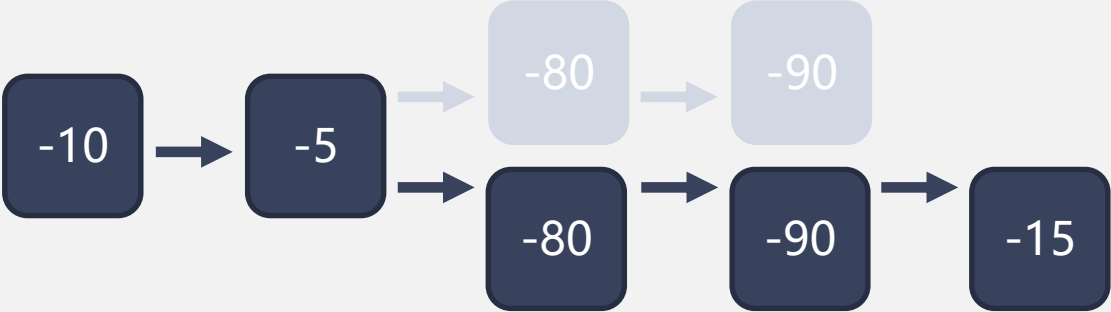
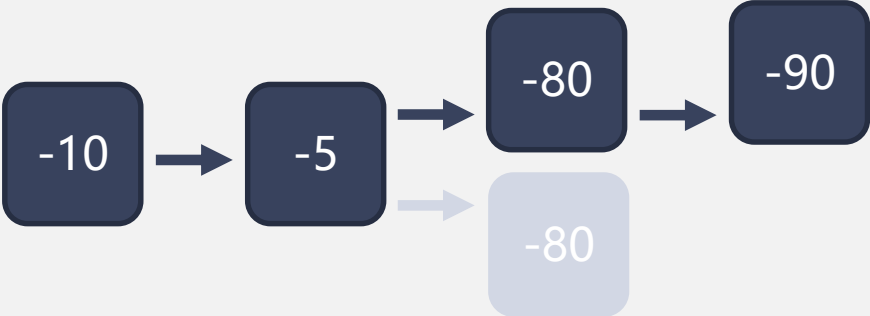




How to prevent double spending for Libra

黃郁傑 林宸熏 楊政道

Double Spending





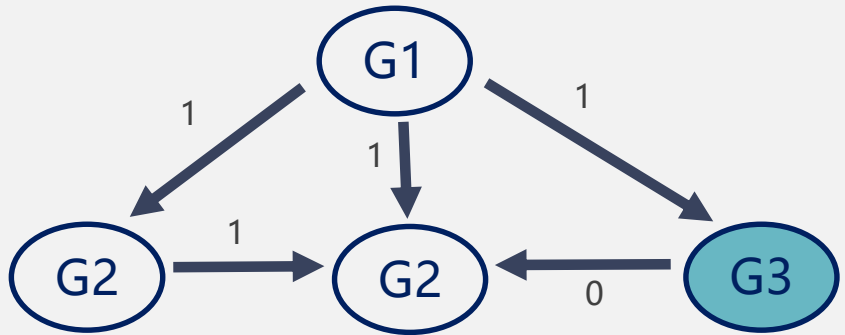
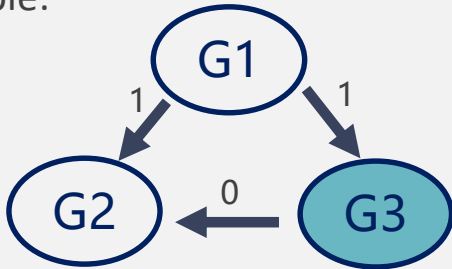
BFT protocol



Byzantine failure

- ◆ Traitor nodes send conflicting messages
 - Which leads to an incorrect result
- ◆ Cause:
 - Flaky nodes
 - Malicious nodes

- ◆ Example:



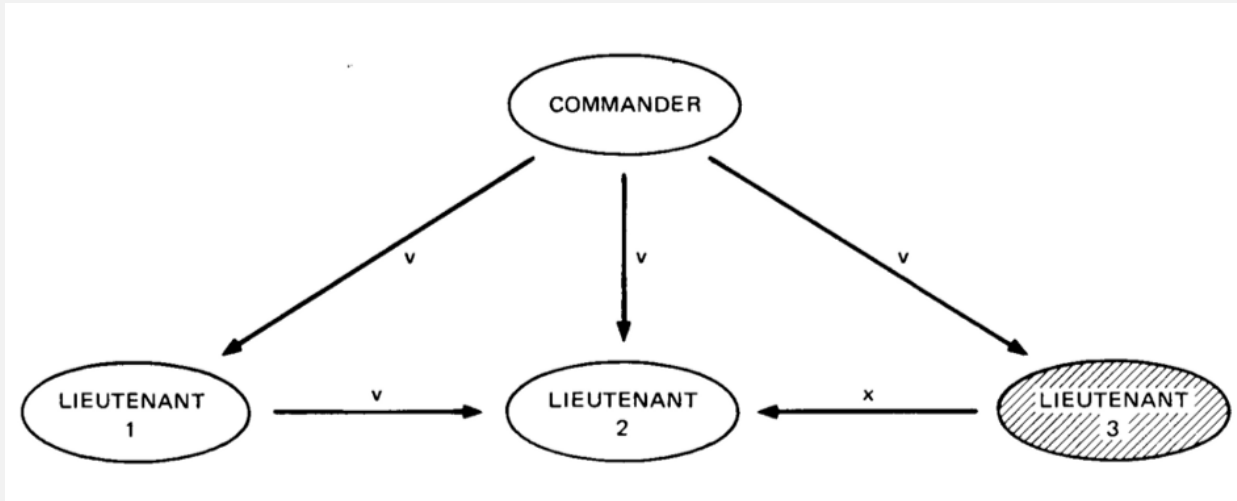
The question



If we want all the generals to get the “same” value...

How many traitors can we
be tolerated?

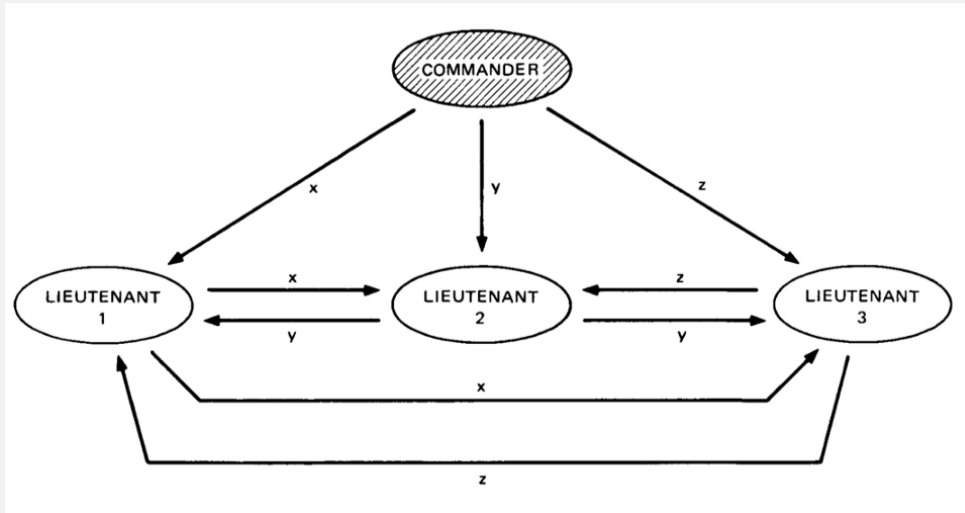
Situation 1 (when commander is not a traitor)



➔ $L2 \leftarrow \max(v, v, x) == v$

➔ L2 can get the “correct” value from commander

Situation 2 (when commander is a traitor)



➔ $L1 \leftarrow \max(x,y,z) \mid L2 \leftarrow \max(x,y,z) \mid L3 \leftarrow \max(x,y,z)$

➔ L2 can get the “same” value from commander

Result

Lemma: **No** solution for $3m+1$ generals with $>m$ traitors.

Proof:

1. Assume solution exists.
2. Use solutions to solve **1** traitors **3** generals case.



We know 2 is impossible.



Hence solution must not exist.

Byzantine Fault Tolerance (BFT)

How might you build such a system ??



The Lamport, Pease and Shostak Algorithm



The Lamport, Pease and Shostak Algorithm

Algorithm OM(0)

1. The general sends his value to every lieutenant.
2. Each lieutenant uses the value he receives from the general.

Algorithm OM(m), $m > 0$

1. The general sends his value to each lieutenant.
2. For each i , let v_i be the value lieutenant i receives from the general. Lieutenant i acts as the general in Algorithm OM($m-1$) to send the value v_i to each of the $n-2$ other lieutenants.
3. For each i , and each $j \neq i$, let v_i be the value lieutenant i received from lieutenant j in step 2 (using Algorithm ($m-1$)). Lieutenant i uses the value majority (v_1, v_2, \dots, v_n).<>

Lamport's Algorithm Definition

Running time (time complexity)

m	Time complexity
0	$O(n)$
1	$O(n^2)$
2	$O(n^3)$
3	$O(n^4)$





Practical Byzantine Fault Tolerance (PBFT)



FEATURE OF PRACTICAL BYZANTINE FAULT TOLERANCE



Total nodes > 3 X faulty nodes



Every node know each other



Each action has a sequence number



Only one leader node



Nodes can ask to change leader (view-change)



WHY N MUST BE BIGGER THAN 3F

N : total nodes

F : faulty nodes (may be **faulty** or **not responding**)

01 Threshold = $N - F$

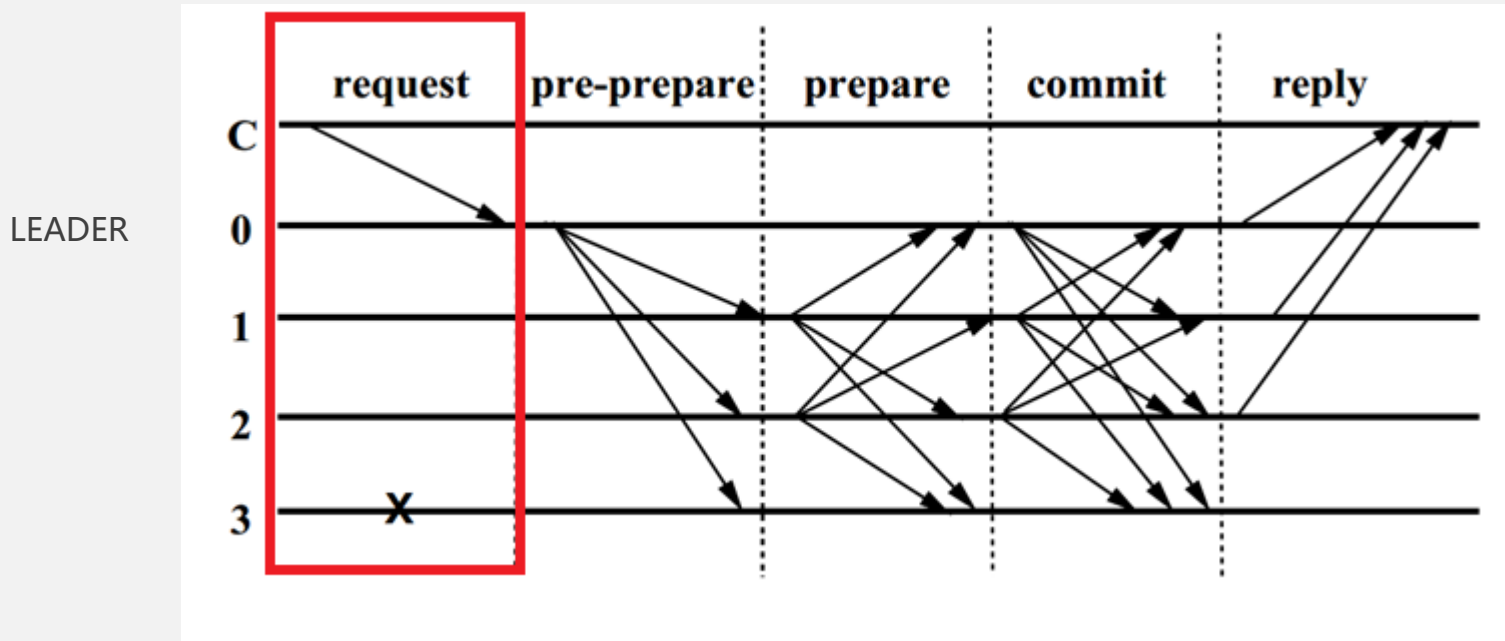
Must make the decision within $N - F$ nodes

02 $(N - F) - F > F$

To ensure normal nodes more than faulty nodes

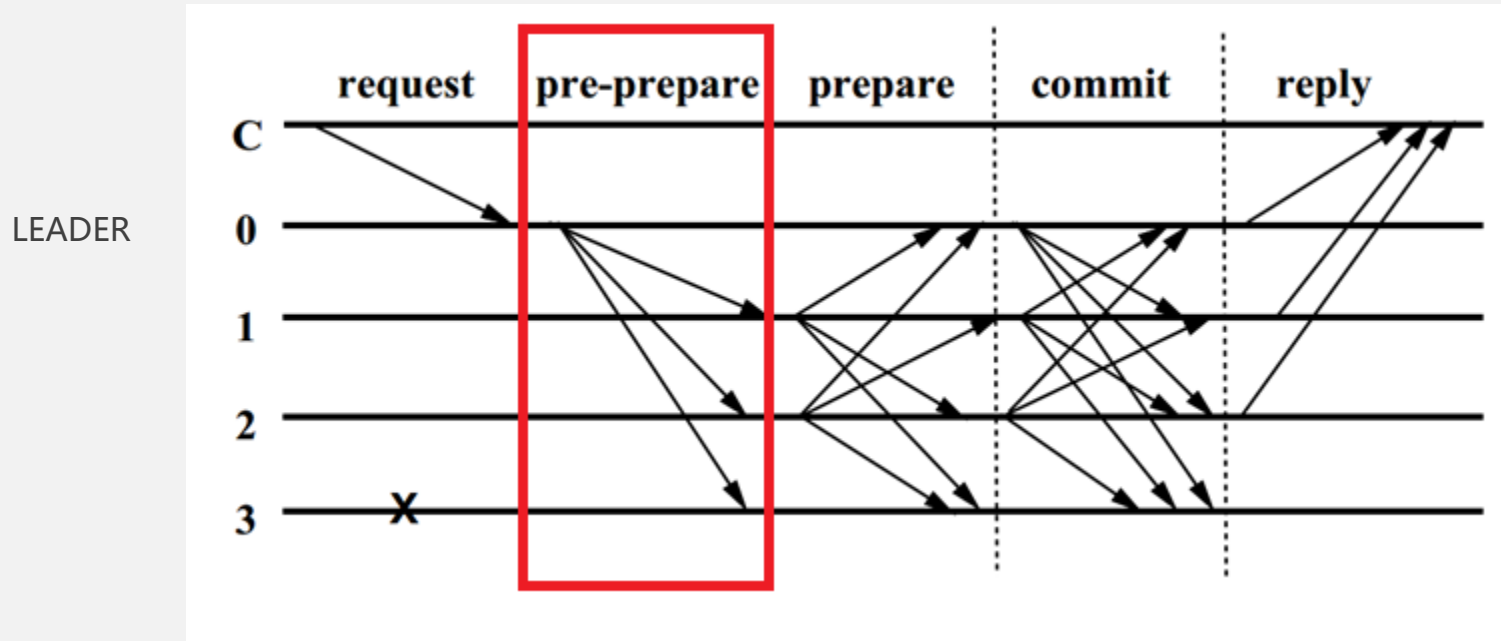
OPERATION OF PRACTICAL BYZANTINE FAULT TOLERANCE

1. Request



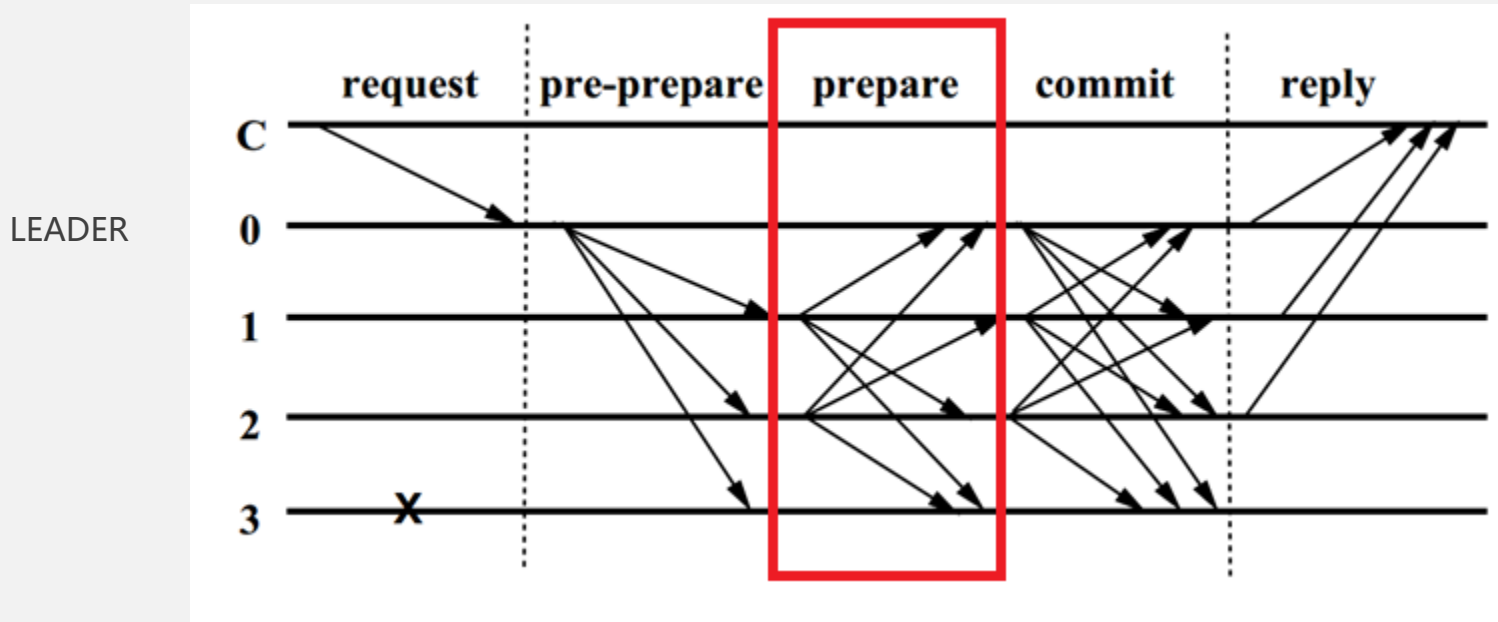
OPERATION OF PRACTICAL BYZANTINE FAULT TOLERANCE

2.Pre-prepare



OPERATION OF PRACTICAL BYZANTINE FAULT TOLERANCE

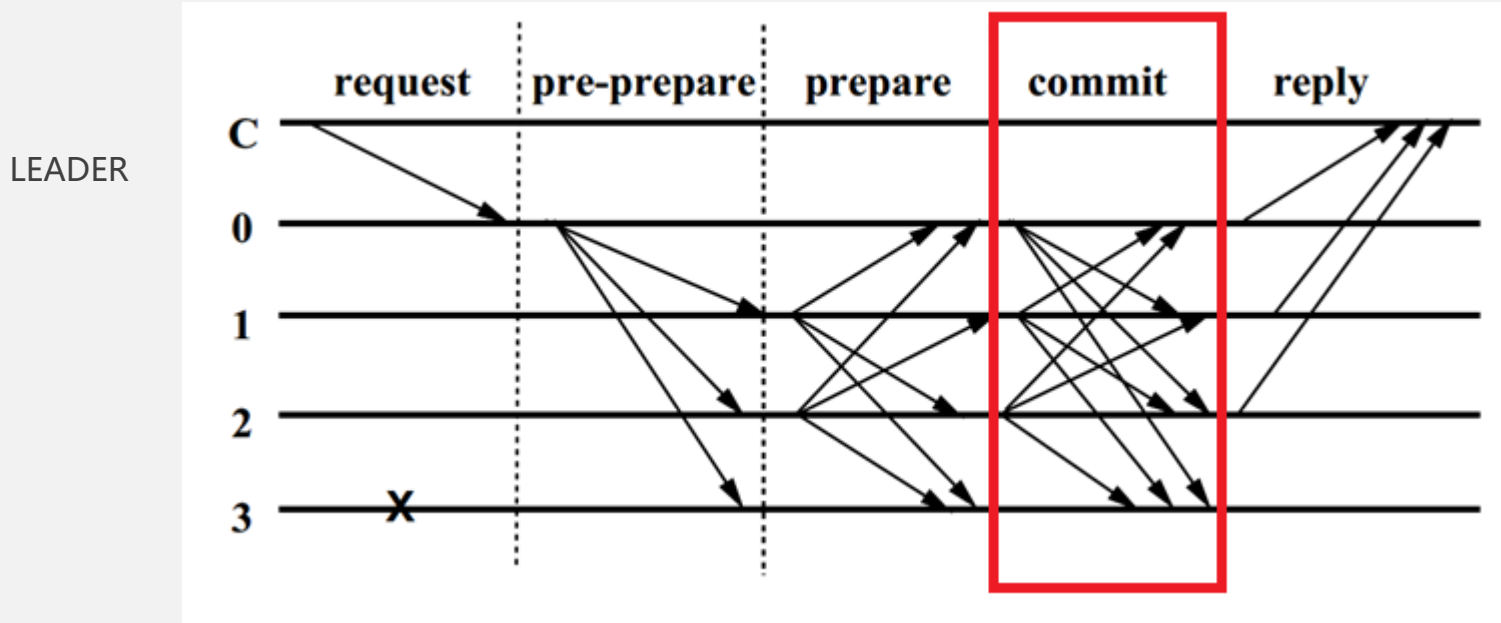
3.Prepare: decide whether the request is legal
one node can make decision if it get more than $2F+1$ votes



OPERATION OF PRACTICAL BYZANTINE FAULT TOLERANCE

4.Commit: vote

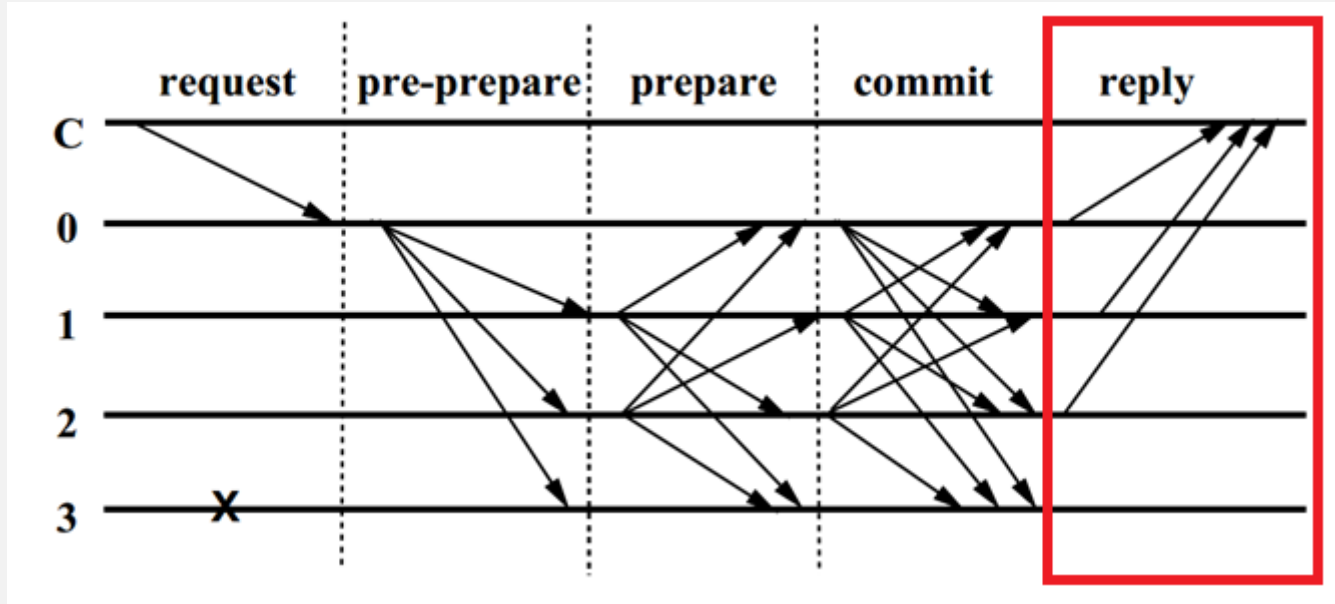
one node can make decision if it get more than $2F+1$ votes



OPERATION OF PRACTICAL BYZANTINE FAULT TOLERANCE

5. Reply: make final decision if it get more than $F+1$ votes

LEADER

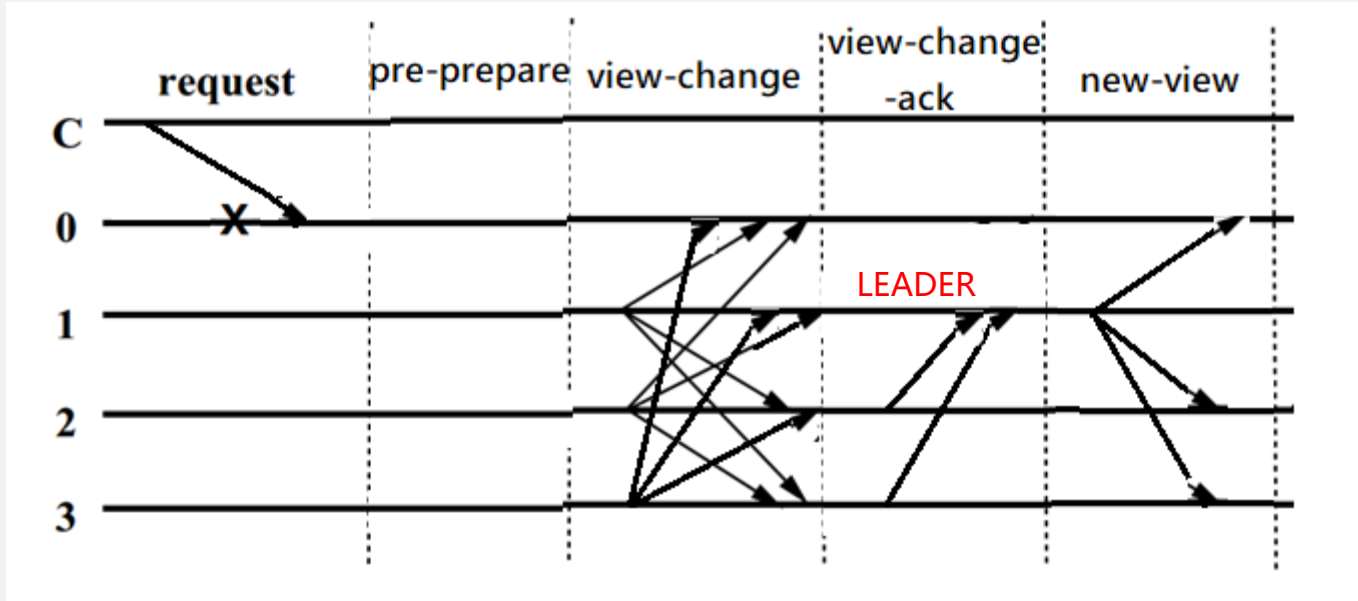


VIEW CHANGE



When 0 is a bad leader

LEADER



ADVANTAGES OF PRACTICAL BYZANTINE FAULT TOLERANCE

01

Contain fault tolerance

02

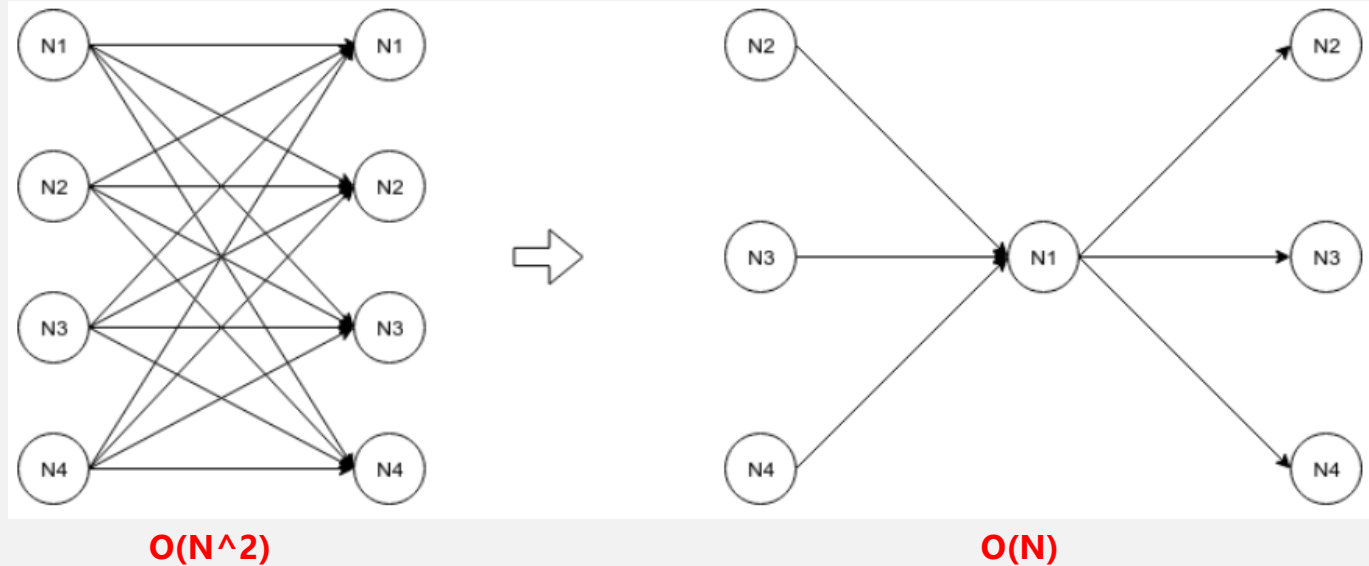
No need to launch token

DISADVANTAGES OF PRACTICAL BYZANTINE FAULT TOLERANCE

01 Leader based

02 Closure property

03 High complexity





CONTENTS

01 Round

What do validators and leader do in a round?

02 Partial synchrony

Synchrony? Asynchrony? WE WANT BOTH!

03 Record

How do Libra build a “chain”?

04 Protocol

Let's take a closer look in Libra!



Round





user

transaction



leader

block



nodes



$N - f$ vote

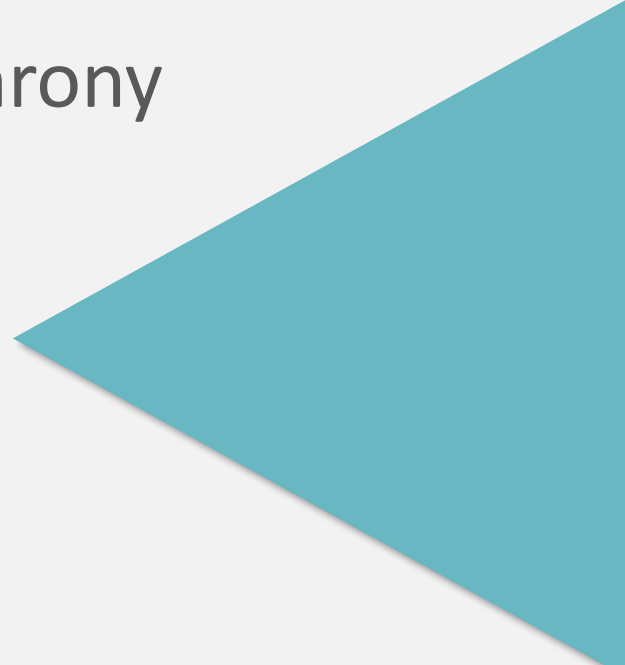


$f+1$ timeout





Partial synchrony



Internet delay

DoS

Communication delay





Asynchrony: waiting for infinite time

BFT: waiting for at most Δ

LBFT: waiting for $\Delta + \max(t, GST)$



Record



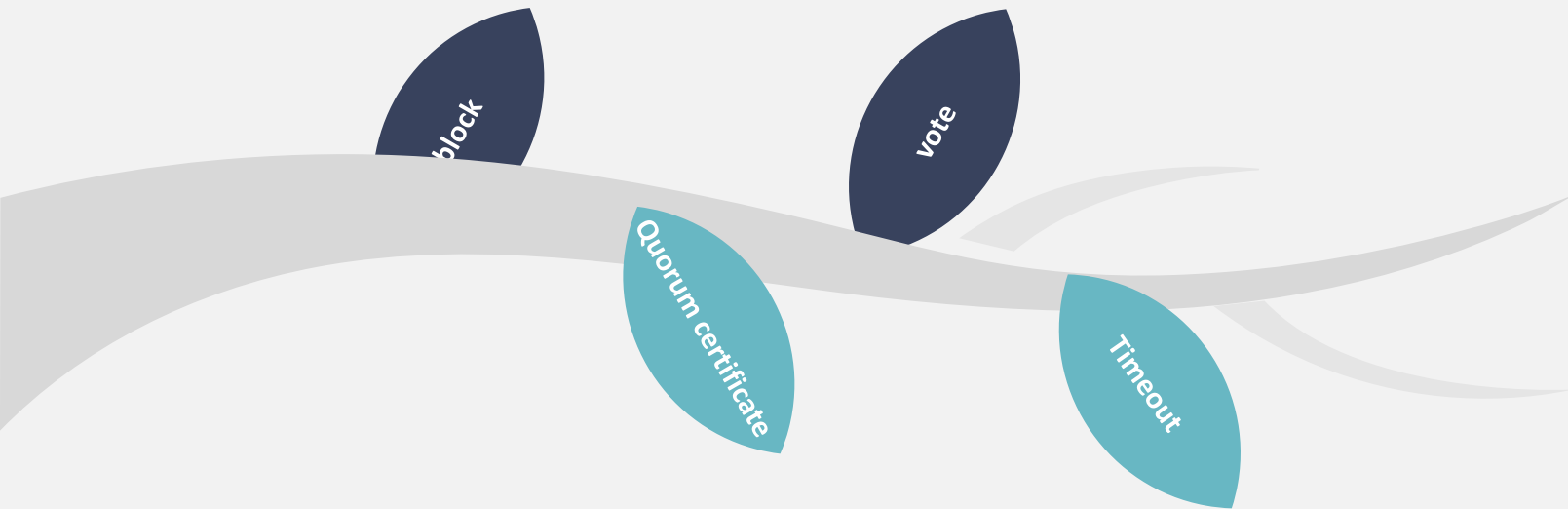
Records

Block

由leader在指定的round提案

vote

Node投票同意block和執行後的state



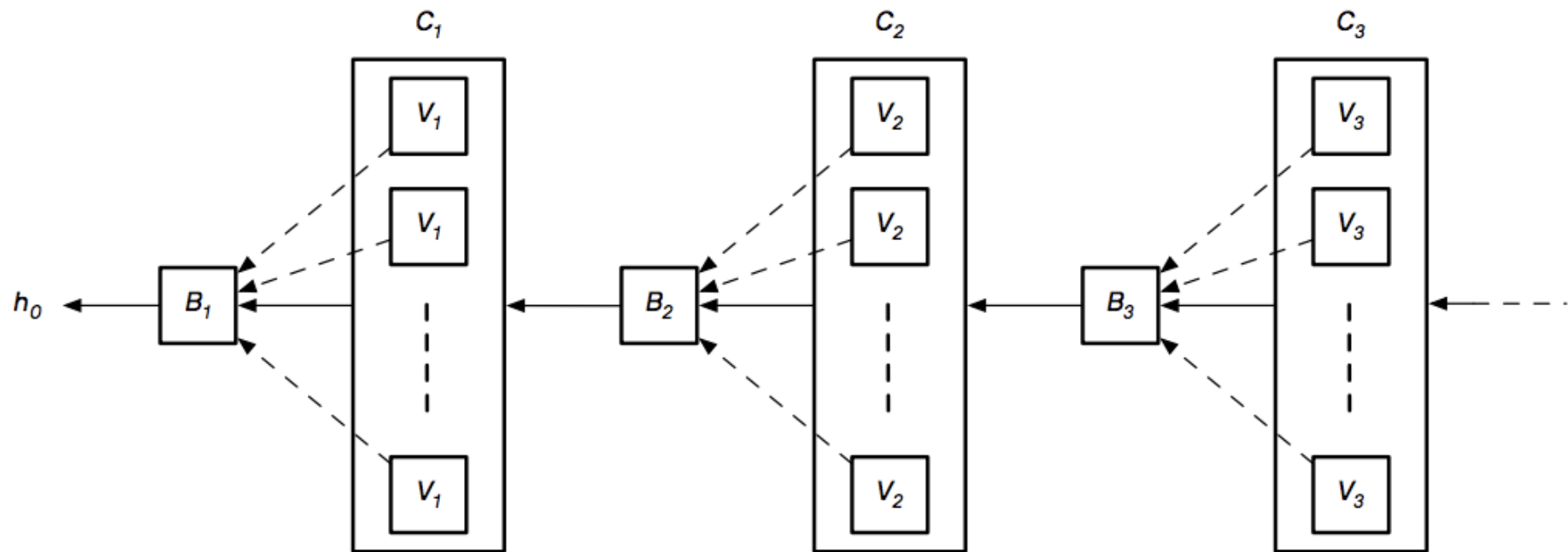
Quorum certificate

Vote的集合並帶有選擇性
欄位commitment

Timeout

投票表示更換leader

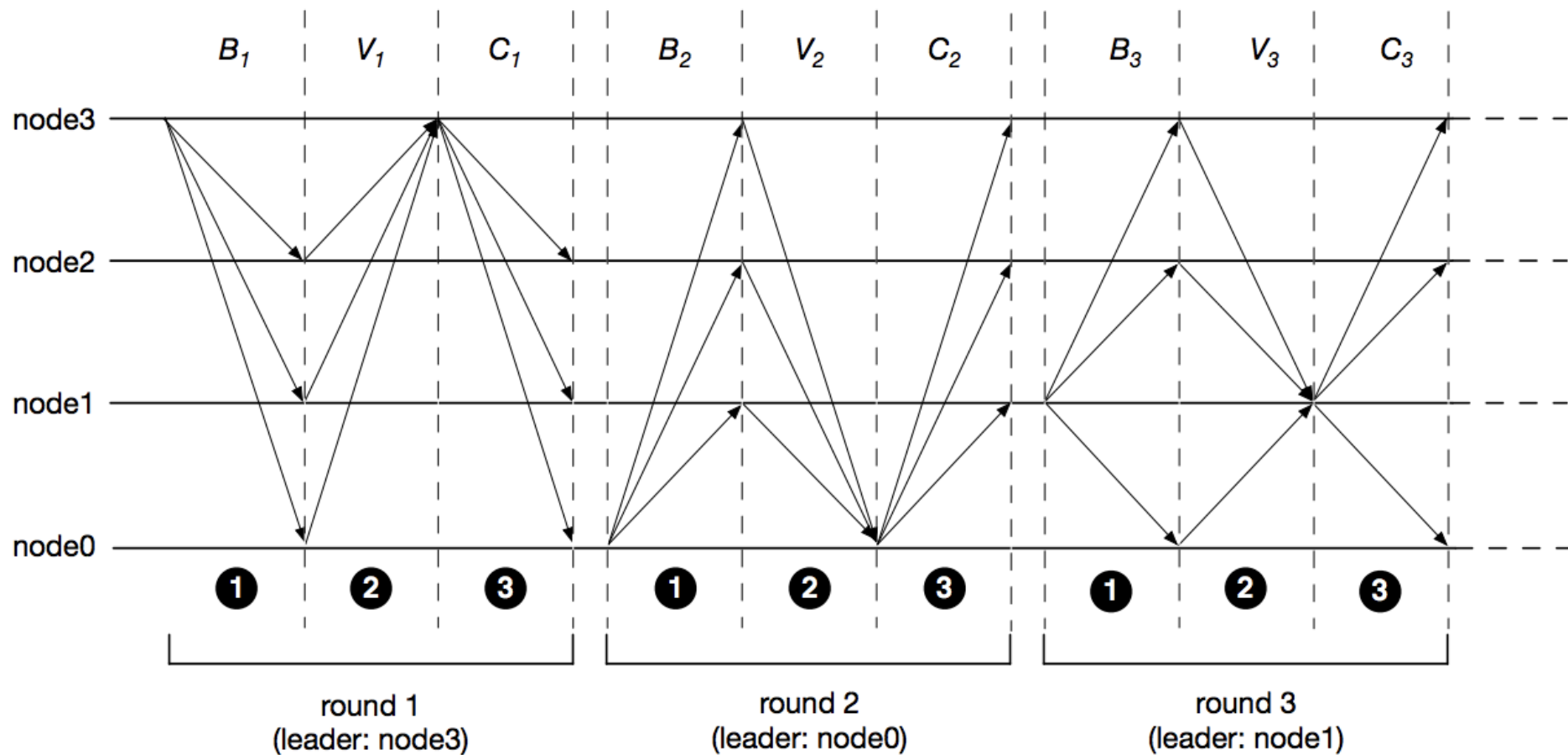
每個部分都有author和signature





Protocol





聖人大盜 電影心得

這部聖人大盜的電影中，多次提到區塊鏈應用，主角團們希望透過區塊鏈的技術結合金融革新，來達到他們創業的目標，同時也是一種追逐夢想的故事。然而看似"去中心化"的中心理念，卻在後續跟政府的合作報告中無法實現，也因為有這些弊端，主角團們才開始思考在夢想與現實妥協是不是一個正確的決定。最後，主角團為了不讓這些社會上游的人持續把持住這個社會，決定最初反撲。我覺得，整片故事中，傳遞了很強的金融體系不公正的問題，因此，在這樣的時空背景之下，區塊鏈變成一種改正社會風氣的正義，也透過主角團們實現區塊鏈的應用來傳遞社會正義的實現。



Thanks for listening!

